

Viewpoints on the p -curvature

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E-functions, **G**-functions and **P**eriods

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Viewpoints on linear differential equations

$$a_r(x) \cdot y^{(r)} + a_{r-1}(x) \cdot y^{(r-1)} + \cdots + a_1(x) \cdot y' + a_0(x) \cdot y = 0$$

Differential operators

$$L = a_r \partial^r + \cdots + a_1 \partial + a_0 \in k(x) \langle \partial \rangle$$

$$L(y) = 0$$

Differential systems

$$Y \in k(x)^r, A \in M_r(k(x))$$

$$Y' + AY = 0$$

Modules with connections

M vector space over $k(x)$

$\partial : M \rightarrow M$ such that

$$\partial(fm) = f'm + f\partial(m)$$

Some examples

$$\Rightarrow y' - y = 0 \qquad y = e^x$$

$$\Rightarrow x \cdot y' - n \cdot y = 0 \qquad y = x^n$$

$$\Rightarrow x(x-1) \cdot y'' + (c - (a+b+1)x) \cdot y' - ab \cdot y = 0 \qquad \sum_{n=0}^{\infty} \frac{(a)_n (b)_n}{(c)_n n!} x^n$$

$$\Rightarrow (x\partial - \alpha_1) \cdots (x\partial - \alpha_n) - x(x\partial - \beta_1) \cdots (x\partial - \beta_n) \sum_{n=0}^{\infty} \frac{(\alpha_1)_n \cdots (\alpha_n)_n}{(\beta_1 + 1)_n \cdots (\beta_n + 1)_n} x^n$$

Theorem (Beukers-Heckmann)

The above differential equation has a basis of algebraic solutions iff:

- \Rightarrow the α_i and the β_j are $2n$ distinct elements of \mathbb{Q}/\mathbb{Z} ,
- \Rightarrow for all r , the sets $A_r = \{r\alpha_1, \dots, r\alpha_n\}$ and $B_r = \{r\beta_1, \dots, r\beta_n\}$ are intertwined in \mathbb{Q}/\mathbb{Z} .

Viewpoints on linear differential equations

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Differential operators

$$L = a_r \partial^r + \cdots + a_1 \partial + a_0 \in k(x) \langle \partial \rangle$$

$$L(y) = 0$$

p -curvature

$$\partial^p \bmod L$$

Differential systems

$$Y \in k(x)^r, A \in M_r(k(x))$$

$$Y' + AY = 0$$

p -curvature

A_p avec :

$$A_1 = A$$

$$A_{i+1} = A'_i + AA_i$$

Modules with connections

M vector space over $k(x)$

$\partial : M \rightarrow M$ such that

$$\partial(fm) = f'm + f\partial(m)$$

p -curvature

$$\partial^p : M \rightarrow M$$

$$\partial^p(fm) = f^{(p)}m + f\partial^p(x) = f\partial^p(m)$$

p -curvature and solutions

Theorem (Cartier)

$$\begin{aligned} & \text{dimension}^{\text{over the subfield of constants}} \text{ of the spaces of solutions} \\ = & \text{dimension}^{\text{over } k(x)} \text{ of the kernel of the } p\text{-curvature} \end{aligned}$$

p -curvature and solutions

Theorem (Cartier)

$$\begin{aligned} & \text{dimension}^{\text{over the subfield of constants}} \text{ of the spaces of rational solutions} \\ = & \text{dimension}^{\text{over } k(x)} \text{ of the kernel of the } p\text{-curvature} \end{aligned}$$

p -curvature and solutions

Theorem (Cartier)

$$\begin{aligned} & \text{dimension}^{\text{over the subfield of constants}} \text{ of the spaces of algebraic solutions} \\ = & \text{dimension}^{\text{over } k(x)} \text{ of the kernel of the } p\text{-curvature} \end{aligned}$$

p -curvature and solutions

Theorem (Cartier)

$$\begin{aligned} & \text{dimension}^{\text{over the subfield of constants}} \text{ of the spaces of series solutions} \\ = & \text{dimension}^{\text{over } k(x)} \text{ of the kernel of the } p\text{-curvature} \end{aligned}$$

p -curvature and solutions

Theorem (Cartier)

$$\begin{aligned} & \text{dimension} \overset{\text{over the subfield of constants}}{\wedge} \text{ of the spaces of series solutions} \\ = & \text{dimension} \overset{\text{over } k(x)}{\wedge} \text{ of the kernel of the } p\text{-curvature} \end{aligned}$$

(M, ∂) module with connection

$$\dim_{k(x^p)} \ker \partial = \dim_{k(x)} \ker \partial^p$$

$$\ker \partial^p = k(x) \otimes_{k(x^p)} \ker \partial$$

$$Y' + AY = 0 \text{ with } Y = Y_0 + Y_1x + Y_2x^2 + \dots$$

$$\text{Recurrence: } nY_n = f_n(Y_0, \dots, Y_{n-1})$$

Necessary condition to the existence of solutions:

$$f_p(Y_0, \dots, Y_{p-1}) = 0$$

Grothendieck conjecture

Let $L \in \mathbb{Q}(x)\langle\partial\rangle$.

The differential equation $L(y) = 0$
has a basis of algebraic solutions

iff the p -curvature of $L \bmod p$
vanishes for almost all prime p

$$\partial^p \equiv 0 \pmod{L, p}$$

Let $L \in \mathbb{Q}[X]$ separable.

The polynomial L splits over \mathbb{Q}

iff the polynomial $L \bmod p$
splits over \mathbb{F}_p for almost all prime p

$$X^p \equiv X \pmod{L, p}$$

Katz Conjecture

The p -curvatures
are dense in
the Lie algebra
of the Galois group

Chebotarev

Every element of
the Galois group
is a Frobenius
at some p

Naive computation of the p -curvature

Reminder

The p -curvature of $Y' = AY$ is A_p with

$$A_1 = A$$

$$A_{i+1} = A'_i + AA_i$$

First order

The p -curvature of $y' = a(x)y$ is

$$-\frac{d^{p-1}a(x)}{d^{p-1}x} - a(x)^p$$

In general

The recurrence gives an algorithm with complexity $\tilde{O}(dr^\omega p^2)$

To be compared with the size of A_p which is $O(dr^2 p)$

Fast computation of the p -curvature

$$Y' + AY = 0 \text{ with } Y = Y_0 + Y_1x + Y_2x^2 + \dots$$

$$\text{Recurrence: } nY_n = f_n(Y_0, \dots, Y_{n-1})$$

Necessary condition to the existence of solutions:

$$f_p(Y_0, \dots, Y_{p-1}) = 0$$

$$k[[x]]^{\text{dp}} = \left\{ f(x) = \sum_{i=0}^{\infty} c_i \frac{x^i}{i!}, \ c_i \in k \right\}$$

$$x^p = 0 \text{ in } k[[x]]^{\text{dp}}$$

Proposition

Let $A \in M_r(k(x))$ with no pole at 0.

The differential system $Y' = AY$ admits
a fundamental system of solutions over $k[[x]]^{\text{dp}}$

Fast computation of the p -curvature (cont.)

Observation

Let $A \in M_r(k(x))$ with no pole at 0.

For all solution $Y = Y_0 + Y_1x + Y_2\frac{x^2}{2!} + \cdots + Y_p\frac{x^p}{p!} + \cdots$
of $Y' = AY$,

we have $A_p Y_0 = -Y_p$ in $k[x]/x^p$

Corollary (Bostan, C., Schost)

There exists an algorithm that compute the p -curvature
in complexity $\tilde{O}(dr^\omega p)$

There exists an algorithm that compute the
the similarity invariant of the p -curvature
in complexity $\tilde{O}(d^{\omega+\frac{3}{2}}r^{\omega+1}\sqrt{p})$

Method

Use the theorem to compute $A_p \bmod (x-a)^p$ for enough values of a

The characteristic polynomial of the p -curvature

We introduce the *Euler operator* $\theta = x\partial$

$$k[x]\langle\partial\rangle \longrightarrow k[\theta]\langle x^{\pm 1}\rangle \quad [x\theta = (\theta-1)x]$$

$$\partial \longmapsto x^{-1}\theta$$

$$L \longmapsto R$$

$$A \in M_r(k(x)) \\ \rightsquigarrow A_p$$

$$B \in M_s(k(\theta)) \\ \rightsquigarrow B_p = B(\theta) B(\theta+1) \cdots B(\theta+p-1)$$

Renormalized characteristic polynomials

$$P(T, x^p) = \det(T - A_p) \cdot a_r(x)^p$$

$$Q(T, \theta^p - \theta) = \det(T - B_p) \cdot b_s(\theta) b_s(\theta+1) \cdots b_s(\theta+p-1)$$

Theorem (Bostan, C., Schost)

We have $P(T, U) = Q(U, TU)$

There exists an algorithm that computes $\det(T - A_p)$
in complexity $\tilde{O}(d(d+r)^\omega \sqrt{p})$

Example: hypergeometric operators

$$\begin{aligned}\mathcal{H}(\alpha_1, \dots, \alpha_n; \beta_1, \dots, \beta_n) \\&= (x\partial - \alpha_1) \cdots (x\partial - \alpha_n) - x(x\partial - \beta_1) \cdots (x\partial - \beta_n) \\&= (\theta - \alpha_1) \cdots (\theta - \alpha_n) - x \cdot (\theta - \beta_1) \cdots (\theta - \beta_n)\end{aligned}$$

More generally, consider $\mathcal{H}(\alpha; \beta) = \alpha(\theta) - x \cdot \beta(\theta)$

$$B = \left(\frac{\alpha(\theta)}{\beta(\theta)} \right) \quad ; \quad B_p = \left(\frac{\alpha(\theta) \alpha(\theta+1) \cdots \alpha(\theta+p-1)}{\beta(\theta) \beta(\theta+1) \cdots \beta(\theta+p-1)} \right)$$

$$\begin{aligned}Q(T, \theta^p - \theta) \\&= \det(T - B_p) \cdot \beta(\theta) \beta(\theta+1) \cdots \beta(\theta+p-1) \\&= \beta(\theta) \beta(\theta+1) \cdots \beta(\theta+p-1) \cdot T \\&\quad - \alpha(\theta) \alpha(\theta+1) \cdots \alpha(\theta+p-1)\end{aligned}$$

We retrieve $P(T, x^p) \propto \det(T - A_p)$
after the changes of variables $T \rightsquigarrow x^p$ and $\theta^p - \theta \rightsquigarrow x^p T$

Merci **pour** **votre** **attention**